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The proper treatment of linguistic ambiguity in ordinary algebra

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(1) *bank* (‘financial institute’, ‘strip of land along a river’)

Old puzzle:

- If we treat ambiguity as a semantic phenomenon (i.e. assume the existence of *ambiguous meanings*),
- the question arises: what is the meaning of ambiguity?
- In particular, one has to account for the property of *universal distribution* (see below)

Answers:

- We provide a simple axiomatization which captures all combinatorial aspects of ambiguity in the context of Boolean algebras.
- But: from these axioms (which are correct beyond doubt) a lot of properties follow which are *not appropriate* for ambiguity.

New puzzle:

- If the most obvious properties of ambiguity lead to properties which are obviously wrong – how do we get out of this?
- One possibility: ambiguity is not *total*, rather a partial operator.
- Another possibility: natural language semantics is not Boolean in nature.

- 1 Linguistic ambiguity
- 2 The semantics of linguistic ambiguity
- 3 Ambiguous algebras
- 4 Structure theory I: Uniformity
- 5 The meaning of uniformity
- 6 Structure theory II: Completions
- 7 Conclusion and further work

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Linguistic ambiguity

Exponents of natural language give rise to two or more sharply distinguished meanings.

- (2) a. *bank* ('financial institute', 'strip of land along a river')
 b. *kick the bucket* ('kick the bucket', 'die')
 c. *see the man with the telescope* ('see with', 'man with')
 d. *every boy loves a movie* ($\forall\exists$, $\exists\forall$)

Syntactic versus semantic approach

Say:

e	is an EXPONENT	(<i>bank</i>)
m_1	is one MEANING of e	(‘financial institute’)
m_2	is another MEANING of e	(‘river bank’)
(e, m)	is a SYMBOL	

Syntactic approach

Two separate symbols (e, m_1) and (e, m_2) .

Semantic approach

One ambiguous symbol $(e, m_1 \parallel m_2)$.

Question: What does \parallel actually mean?

Advantages of the semantic approach

Semantic approach

One ambiguous symbol ($e, m_1 \parallel m_2$).

Advantages:

- function from form to meaning
- more succinct lexicon
- interaction of ambiguity, meaning composition, and inference

(3) a. *The federal agency decided to take the project under its well-muscled, federally-funded wing.*^[1]

b. *We pulled his cross-gartered leg.*^[1]

(4) *The first thing that strikes a stranger in New York is a big car.*^[3]

What \parallel doesn't mean

Question: What does \parallel actually mean?

Quick shot: Disjunction!

$$m_1 \parallel m_2 \equiv m_1 \vee m_2$$

$$\text{Hence: } \neg(m_1 \parallel m_2) \equiv \neg(m_1 \vee m_2) \equiv \neg m_1 \wedge \neg m_2$$

$$\text{Yet intuitively: } \neg(m_1 \parallel m_2) \equiv \neg m_1 \parallel \neg m_2$$

(5) *#There is no bank.*

‘There is no financial institute and there is no strip of land along the river.’

(see also Pinkal [2], Poesio [3], and Stallard [4])

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The semantics of linguistic ambiguity

We investigate the meaning of \parallel in the context of **Boolean algebras**.

- This is an important restriction: Boolean algebras correspond (via algebraization) to classical logic – which might not necessarily be the right logic for NL-semantics.
- Still, the vast majority of semanticists interpret connectives in classical logic – so for all choices, this seems to be the most natural!
- In the algebraic approach, algebraic \leq (defined by $a \leq b \Leftrightarrow a \wedge b = a$) corresponds to logical \vdash and semantic \models

The semantics of linguistic ambiguity

Question: What are the semantic properties of \parallel ?

In terms of **denotation**, there is a fundamental difference between disjunction and ambiguity: the meaning of an ambiguous statement depends on the underlying (often unknown) **intention** of the speaker:

Intentionality

$$a \parallel b \leq a \vee b \quad (6)$$

$$a \leq a \vee b \quad (7)$$

$$\textbf{But: } a \not\leq a \parallel b \quad (8)$$

I need some money! $\not\leq$ I need some dough!

I need some pastry or some money! \neq I need some dough!

The semantics of linguistic ambiguity

The **combinatorial properties** of ambiguity are different from disjunction: ambiguity has the property of universal distribution:

Universal distribution

$$\sim(a \parallel b) = \sim a \parallel \sim b \quad (9)$$

$$(a \parallel b) \vee c = (a \vee c) \parallel (b \vee c) \quad (10)$$

$$(a \parallel b) \wedge c = (a \wedge c) \parallel (b \wedge c) \quad (11)$$

$$(a \parallel b) \rightarrow c = (a \rightarrow c) \parallel (b \rightarrow c) \quad (12)$$

$$a \rightarrow (b \parallel c) = (a \rightarrow b) \parallel (a \rightarrow c) \quad (13)$$

The semantics of linguistic ambiguity

There are some additional properties of \parallel :

Associativity

$$(a \parallel b) \parallel c = a \parallel (b \parallel c) \quad (\text{ass})$$

Idempotence

$$a \parallel a = a \quad (\text{id})$$

Commutativity (arguable)

$$a \parallel b = b \parallel a \quad (\text{com})$$

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Uniform Usage

The following hypothesis of uniform usage is necessary for an algebraic treatment of ambiguity:

Uniform usage (UU)

In a given context, an ambiguous statement is used consistently in **only one** sense.

This leaves of course many things unspecified (as context), but allows us to treat \parallel as an algebraic operator (which is a function!)

Ambiguous algebra

An **AMBIGUOUS ALGEBRA** is a structure $\mathbf{A} = (A, \wedge, \vee, \sim, \parallel, 0, 1)$, where $(A, \wedge, \vee, \sim, 0, 1)$ is a Boolean algebra, and \parallel is a binary operation for which the following holds:

$$\sim(a \parallel b) = \sim a \parallel \sim b \quad (\parallel 1)$$

$$a \wedge (b \parallel c) = (a \wedge b) \parallel (a \wedge c) \quad (\parallel 2)$$

$$\text{At least one of } a \leq a \parallel b \text{ or } b \leq a \parallel b \text{ holds} \quad (\parallel 3)$$

Ambiguous algebras: a peculiar axiom

At least one of $a \leq a \parallel b$ or $b \leq a \parallel b$ holds $(\parallel 3)$

Note that $(\parallel 3)$ is a **disjunction**!

- This entails, among other, that there is no *free ambiguous algebra*, a central tool in general algebra.
- Put differently, in *every* ambiguous algebra some equalities hold which do not hold in all ambiguous algebras (this nicely models the epistemic aspect of ambiguity)
- To the best of our knowledge, axioms of this kind have not been considered in general algebra so far. Any algebraist know better?

Questions regarding the axiomatization

- 1 Do these axioms entail all properties we find intuitively true for ambiguity?
 - As far as we can see, clearly yes.
- 2 Do they imply some properties we find intuitively incorrect for ambiguity in general?
 - Unfortunately, also clearly yes.
- 3 Do non-trivial algebras exist which satisfy these axioms? (That is, for example, algebras with more than one element?)
 - Clearly yes, but if we add commutativity for \parallel , then no.
- 4 Are there ambiguous algebras, where $a \parallel b \neq a$ and $a \parallel b \neq b$?
 - No, there are not.

Ambiguous algebras: An example

Take the obvious Boolean algebra over the set $\{0, a, b, 1\}$. Put

$$\begin{aligned}a \parallel b &= a; & b \parallel a &= b; & 0 \parallel a &= 0; & 1 \parallel a &= 1; \\a \parallel 1 &= a; & b \parallel 1 &= b; & 0 \parallel b &= 0; & 1 \parallel b &= 1; \\a \parallel 0 &= a; & b \parallel 0 &= b; & 0 \parallel 1 &= 0; & 1 \parallel 0 &= 1;\end{aligned}$$

\wedge -distribution holds:

$$a = a \parallel b = (a \parallel b) \wedge a = a \parallel 0 = a$$

$$0 = 0 \parallel 1 = (0 \parallel 1) \wedge a = 0 \parallel a = 0$$

and so on, same for \vee, \sim

We thus have a proper non-trivial 4-element algebra.

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Structure theory I: Uniformity

Lemma 1

In every ambiguous algebra \mathbf{A} and for all $a, b \in A$,
either $a = a \parallel b$ or $b = a \parallel b$.

No commutativity for \parallel :

assume $a \parallel \sim a = a$

then $\sim a \parallel a = a$

but also $\sim(\sim a \parallel a) = \sim\sim a \parallel \sim a = a \parallel \sim a = a$

hence $\sim a = a$ (which only holds in 1-element algebras)

Linguistically relevant?

(14) *sacré* ('cursed', 'holy')

Structure theory I: Uniformity

Lemma 1

In every ambiguous algebra \mathbf{A} and for all $a, b \in A$,
either $a = a \parallel b$ or $b = a \parallel b$.

Corollary 2

If \mathbf{A} is an ambiguous algebra such that for all $a, b \in A$, $a \parallel b = b \parallel a$,
then A has at most one element.

Corollary 3

For all ambiguous algebras \mathbf{A} , $a, b \in A$, we have

1. $a \parallel a = a$
2. $a \wedge b \leq a \parallel b \leq a \vee b$

Structure theory I: Uniformity

Lemma 4 (Monotonicity of \parallel)

Assume $a' \leq a$ and $b' \leq b$. Then $a' \parallel b' \leq a \parallel b$.

Lemma 7 (Uniformity lemma)

Assume we have an ambiguous algebra \mathbf{A} $a, b \in A$ such that $a \neq b$.

- 1 If $a \parallel b = a$, then for all $c, c' \in A$, we have $c \parallel c' = c$;
- 2 if $a \parallel b = b$, then for all $c, c' \in A$, we have $c \parallel c' = c'$.

Corollary 8

If \mathbf{A} is an ambiguous algebra, $a, b \in \mathbf{A}$, then either for all $a, b \in A$, we have $a \parallel b = a$, or for all $a, b \in A$, we have $a \parallel b = b$.

Hence: Ambiguous algebras are either RIGHT-SIDED or LEFT-SIDED.

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The meaning of monotonicity

Does monotonicity go together with intuitions about ambiguity?

Lemma 4 (Monotonicity of \parallel)

Assume $a \leq a'$ and $b \leq b'$. Then $a \parallel b \leq a' \parallel b'$.

\Rightarrow This roughly means: ambiguity respects logical entailment.

Example 1:

bank means $a \parallel b$

restaurant means c

bank or restaurant means $(a \parallel b) \vee c = (a \vee c) \parallel (b \vee c) = a' \parallel b'$

\Rightarrow If monotonicity were wrong, then *bank* would not (generally) entail *bank or restaurant*!

The meaning of monotonicity

Does monotonicity go together with intuitions about ambiguity?

Lemma 4 (Monotonicity of \parallel)

Assume $a \leq a'$ and $b \leq b'$. Then $a \parallel b \leq a' \parallel b'$.

⇒ This roughly means: ambiguity respects logical entailment.

Example 2:

bank means $a \parallel b = a$

bank or restaurant means $(a \parallel b) \vee c = (a \vee c) \parallel (b \vee c) = a' \parallel b'$

kank means $a' \parallel b' = a \vee c$

⇒ Hence monotonicity is like uniform usage for expressions which are connected by the relation of entailment.

The meaning of uniformity

Does uniformity go together with intuitions about ambiguity?

Lemma 7 (Uniformity lemma, informal)

Ambiguous algebras are **either** left-sided **or** right-sided.

Hence: $a \parallel b = a$ **or** $a \parallel b = b$ for any a, b

Example: *bank* means ‘financial institute’, and **accordingly**
dough means ‘pastry’.

This is completely unintuitive, and a consequence of two things:

- 1 the strong axioms of Boolean algebras, in particular the equality $\sim\sim a = a$, and
- 2 the fact that \parallel is a total operator (i.e., for all a, b , we have an object $a \parallel b$).

Ways out of the dilemma

As there can be hardly any doubt about (\parallel 1–3), there are only two ways out:

- 1 Boolean algebras/classical connectives are inadequate for NL semantics. Most of the results presented here do no longer hold with Heyting algebras/intuitionistic logic.
- 2 Ambiguity is a **partial operator**, and partially ambiguous algebras are the smallest Boolean algebras generated by some ambiguous terms. In particular, the uniformity lemma depends on the existence of the object $0 \parallel 1$, which has no linguistic motivation and need not exist in partial algebras.

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So far, we have exposed restrictions on ambiguous algebras.

In this part, we show results on the existence of (non-trivial) ambiguous algebras.

Structure theory II: Completions

Let id be the identity map. Every Boolean algebra can be completed to an ambiguous algebra without collapsing elements:

Lemma 17 (Completion lemma)

Every Boolean algebra $\mathbf{B} = (B, \wedge, \vee, \sim, 0, 1)$ can be completed to an ambiguous algebra $\mathbf{A} = (B, \wedge, \vee, \sim, \parallel, 0, 1)$ such that the map $\text{id} : \mathbf{A} \rightarrow \mathbf{B}$ is a Boolean algebra isomorphism.

Together with the uniformity lemma, this proves the following:

Corollary 18

For every Boolean algebra $\mathbf{B} = (B, \wedge, \vee, \sim, 0, 1)$, there are exactly two ambiguous algebras \mathbf{A} such that the map $\text{id} : \mathbf{A} \rightarrow \mathbf{B}$ is a Boolean algebra isomorphism.

That is: for every Boolean algebra there is one left-sided and one right-sided completion.

This means that decidability results for Boolean algebras can be transferred to ambiguous algebras.

- every AA-term t , knowing it is interpreted in a left-(right-)sided ambiguous algebra, can be reduced to a Boolean algebra term.
- To check whether $t_1 = t_2$ holds in all ambiguous algebras, just check:
 - 1 for t_1, t_2 terms in a left-sided algebra, for their Boolean reductions $b_l(t_1), b_l(t_2)$, $b_l(t_1) = b_l(t_2)$ holds in all Boolean algebras (NP complete).
 - 2 same for right-sided algebras.
- The equation holds in both previous cases if and only if it holds in all ambiguous algebras.

So what does this mean?

- We see that ambiguous Boolean algebras are quite well-behaved,
- but on the downside, we also see that there is little of interest to say about them, as they are very similar to Boolean algebras.
- This in itself is however a non-trivial result, which probably does not extend to other (weaker) types of algebras.

To cite the paper: “Still we consider it important to have established these results, which are really **[not]** obvious or trivial.”


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Ambiguous algebra

Boolean algebra + \parallel operator and axiomatization

Does it match linguistic intuition?

⇒ **monotonicity lemma:** words in entailment relations are used consistently meaningwise. 

⇒ **uniformity lemma:** ambiguous terms are either right- or left-ambiguous. 

How to make it fit better?

- more general class of algebras? (e.g., Heyting algebras or distributive, modular, or residuated lattices)
- partial \parallel ?

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